# An Improved Discrete Level of Detail Model Through an Incremental Representation 

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#### Abstract

Real-time applications such as computer and video games, virtual reality and scientific simulation require rendering of complex models for realism. Graphics rendering engines include multiresolution modelling techniques to accelerate the visualization process. The Discrete Level of Detail framework (DLoD) is usually the most popular while the Continuous Level of Detail framework (CLoD) is still not as widely used by software developers. In this paper, we first discuss the benefits and drawbacks of both frameworks. Then, we present a model based on coding a discrete number of levels of detail (LoDs), with more LoDs coded than is usual in DLoD, and with an incremental representation, which is often used in CLoD. This model obtains a performance similar to DLoD by providing optimized LoDs for efficient visualization, while the popping effect is imperceptible. We present specific proposals for each of the three main stages involved in multiresolution processing: geometry simplification, construction of the incremental representation and retrieval of either uniform or view-dependent LoDs.

Categories and Subject Descriptors (according to ACM CCS): I.3.5 [Computer Graphics]: Computational Geometry and Object Modeling


## 1. Introduction

Multiresolution modelling has been a matter of growing interest in the last decade. One of its main goals is to accelerate the visualization process [LRC*03]. Multiresolution modelling allows for adjusting the level of detail (LoD) of the scene so maintaining a constant frame rate and assuring interactivity to the final user. Nowadays, real-time applications such as computer and video games, virtual reality or scientific simulation are among the most demanding and tightly optimized graphics applications. As these applications require rendering of complex models for realism, graphics rendering engines include multiresolution modelling techniques, which have become widely used.

The multiresolution modelling techniques presented in the literature are classified under two main frameworks for managing level of detail: Discrete LoD (DLoD) and Continuous LoD (CLoD). DLoD is the most widely used. This frame-
work manages a small number of independent levels of detail (LoDs), where each approximation or LoD represents the original object using a different number of faces. CLoD is introduced as an alternative which provides a wide range (virtually a continuous range [PS97]) of different approximations, such that the LoD can be adapted to the application requirements with a high degree of accuracy. CLoD has been extended to provide view-dependent LoDs, which is sometimes considered as a third framework [LRC*03].

In order to introduce our proposal we first compare DLoD and CLoD to understand their benefits and drawbacks. We have arranged a series of processing stages (figure 1):

## Off-line process

- Simplification. The object is processed using a simplification tool. In case of DLoD, this process gives the LoDs separately. In case of CLoD, this process gives the se-


Figure 1: Order of operations.
quence of removals necessary to construct a multiresolution representation that provides a wide range of LoDs.

- Construction.

DLoD. This stage consists of only processing each LoD for maximum efficiency in the rendering process by using features depending on the particular hardware targeted, i.e. triangle strips. After completion, standard file formats can be used to store data for each LoD. Consequently, LoDs can be easily used in a wide range of graphic engines.
CLoD. First, a process constructs the multiresolution representation according to the CLoD model to be used. Second, data is processed to take advantage of hardware rendering features. After completion, data is stored in a proper file format, which is generally non standard.

## At run-time

- Load.

DLoD. LoDs are stored in main memory through commonly used data structures and they are easily compiled into hardware command streams as static buffer objects.
CLoD. This framework requires special data structures to arrange data in such a way that LoDs can be retrieved as efficient as possible. These structures differ depending on the particular CLoD model. As in DLoD, data are compiled into hardware command streams.

- LoD Selection. An algorithm selects the most appropriate level of detail to display. This stage is analogous for both frameworks.
- LoD Recovery.

DLoD. As the LoD has been optimized in the off-line process and compiled for efficient rendering in the load stage, no more process is required. However, view-dependent LoDs cannot be retrieved.
CLoD. An algorithm traverses the data structures to recover either uniform or view-dependent LoDs. In order to use the current hardware capabilities an optimization process must be done, thus incurring in an overload, which is even higher in case of view-dependent LoDs.

- Rendering.

DLoD. This framework displays more or less detail than actually required, and presents visual discontinuities when changing between LoDs, known as the popping effect. However, data is provided to GPU highly optimized and maximum performance is obtained.
CLoD. This framework displays the level of detail just required. Popping effect is not perceptible by the user as the difference between consecutive LoDs is very small. However, maximum performance is hard to achieve because LoD optimization is computed at runtime and it is not as good as off-line optimization.

We can conclude that DLoD presents drawbacks as the popping effect and the impossibility of providing view dependent LoDs, and CLoD requires a higher effort in almost every stage and does not reach the performance of DLoD. Furthermore, as graphics hardware evolve, changes or advent of new features require minor changes in DLoD, but major changes, even redesign, in CLoD.

In this paper, we aim to obtain a performance similar to DLoD by providing optimized LoDs for efficient rendering, while the popping effect is imperceptible. We present an improved DLoD (iDLoD) based on the following features: (1) to increase the number of LoDs, such that it depends on the fame rate of the application (which is a constant) and it is independent on the complexity of the object; (2) to encode LoDs using an incremental representation to reduce the storage cost, so that consecutive LoDs share a high number of faces; and (3) the difference between two consecutive LoDs is forced to be a surface patch (see figure 2 ), so that it can be optimized off-line for rendering.

Consequently, a LoD is a reduced set of surface patches, each of which is optimized off-line. At the load stage, data are compiled into static hardware command streams. The complexity of the algorithm for LoD recovery depends on the number of LoDs, which is a constant instead of depending on the size of the object. The provided LoDs fit the required level of detail better than DLoD, thus avoiding the popping effect. Furthermore, in spite of the low granularity, iDLoD can provide view-dependent LoDs.


Figure 2: The left image shows in red the surface patch (61 triangles) selected for simplification at the current LoD. The right image shows the LoD, which is generated by substituting the red triangles by a simplified surface ( 24 triangles). There is a difference of 37 triangles between these two consecutive LoDs.

## 2. Previous work

The original DLoD was presented by Clark in 1976 [Cla76]. This technique began to be used with the main aim of increasing the performance of the graphic system and this was accomplished in applications such as walkthroughs in virtual environments. Recent work has addressed the problem of the popping effect performing smooth transitions between them by means of geo-morphing or blending [SG03, BGB*05]. Numerous methods for CLoD have been presented in the last years. For an overview on the various schemes proposed see [Gar99] [RLB* 02].

El-Sana et al. [ESAV99] presented the Skip Strips model, which maintains a data structure to store the strips that avoids the need to calculate them in real-time. The MTS model [BRR* 04] uses triangle strips as both the storage and the visualization primitive. It consists of a set of multiresolution strips, each of which represents a triangle strip and all its levels of detail. LodStrips model [RCRG06] is entirely based on optimized hardware primitives, triangle strips, and deals with the apparition of degenerate triangles by applying pre-calculated filters.

Ji et al. [JWLL05] suggest a method to select and display several LoDs by using the GPU. In particular, they encode the geometry in a quadtree based on a LoD atlas texture. Masking Strips [ $\left.\mathrm{RCG}^{*} 09\right]$ uses a cache-aware stripification technique to diminish bus traffic and also to apply LoD update routines which remove all the unnecessary degenerate triangles.

Many of the GPU-based continuous models are aimed at view-dependent rendering of massive models. Researchers have recently proposed methods for moving the granularity of the representation from triangles to triangle patches in order to offer view-dependent capabilities for rendering out-of-core models [CGG* 04, SM05]. With a similar objective but with a further GPU exploitation, the GoLD method [ $\mathrm{BGB}^{*} 05$ ] introduces a hierarchy of geometric patches for very detailed meshes with high resolution textures. Recently, Hu et al. [HSH09] presented a fully-GPU implementation of Progressive Meshes [Hop97].

## 3. Heuristic determination of the number of LoDs

In the context of real-time applications, an object that is far away from the viewer will be represented by the lowest level of detail. As the distance decreases the level of detail increases. Despite thousands of LoDs are available, the application usually requires one LoD per frame at most. Let us consider a frame rate of 50 images per second. If a distant object gets very near of the viewer one second later, for instance, only 50 LoDs would be displayed. In case the distant object gets slowly near to the viewer, ten seconds later for example, the application would require 500 LoDs at most when using the same frame rate. However, we can suspect in these cases that the user is not focusing his attention on this object. This perceptual factor is often used to extend the life of a LoD over a number of frames [LRC*03]. In case the application decides to maintain the LoD during 5 frames, only 100 LoDs would be displayed, just two times the frame rate. In case of maintaining the LoD during only 2 frames, 250 LoDs would be enough, just 5 times the frame rate. In any case, the amount of displayed LoDs is very far from the granularity that CLoD provides.

Therefore, we propose the number of LoDs to be a multiple of the frame rate. Experiments have been carried out using a factor of 2 , and three different frame rates: 25,50 and 70 images per second. Thus each object stores 50, 100 and 140 different LoDs.

## 4. Simplification

A simplification process converts a polygonal surface into another surface with a smaller number of polygons. As in our multiresolution model we need a sequence of simplified surface patches, methods based on vertex clustering or merging regions are well suited [Gar99]. However, we used the method proposed by Garland and Heckbert [GH97], which is based on edge collapse, plus a merging process to form surface patches. There is a public domain implementation of the simplification method of Garland and Heckbert so-called Qslim [Gar04], and we use it due to its speed and the quality of the generated meshes.

Our proposal is based on regions that are created and merged until they reach a given simplification error, $E$. For each edge collapse in the sequence provided by $Q$ slim, a new region is created and checked to be merged with any of the previously created regions. Each region is assigned a simplification error, which is the sum of the simplification errors of all the edge collapses grouped for that region. When the simplification error of a region reaches $E$, the region is not further considered for merging.

Let us consider an edge collapse as the removal and creation of triangles in the polygonal surface. For example, the collapse of the edge between triangles $a$ and $j$ in figure 3(a) and 3 (b) would remove triangles $\{a, b, c, d, i, j\}$ and would


Figure 3: Creating and merging regions. First, a region $R_{i}$ is created, (a) $T_{i}^{d}=\{a, b, c, d, i, j\},(b) T_{i}^{c}=\{k, l, m, n\}$. Second, $a$ new region $R_{j}$ is created, (c) $T_{j}^{d}=\{e, f, g, h, m, n\}$, (d) $T_{j}^{c}=\{o, p, q, r\}$. As $I=\left\{T_{i}^{c} \cap T_{j}^{d}\right\}=\{m, n\}$, then $R_{k}=\operatorname{merge}\left(R_{i}, R_{j}\right)$, (e) $T_{k}^{d}=\left\{T_{i}^{d} \cup T_{j}^{d}\right\}-I=\{a, b, c, d, e, f, g, h, i, j\}$, (f) $T_{k}^{c}=\left\{T_{i}^{c} \cup T_{j}^{c}\right\}-I=\{k, l, o, p, q, r\}$.
create triangles $\{k, l, m, n\}$. For each edge collapse, we define a region, $R_{i}$, as the tuple $\left\{T_{i}^{d}, T_{i}^{c}\right\}$ where $T_{i}^{d}$ is the set of removed triangles (figure 3(a)) and $T_{i}^{c}$ is the set of created triangles (figure 3(b)). Once the new region is created, we check the previously created regions to find those regions suitable to be merged with it. So, given two regions, $R_{i}$ and $R_{j}$, where region $R_{i}$ was previously created and $R_{j}$ is a new region, we define the check function as:

$$
I=\operatorname{check}\left(R_{i}, R_{j}\right)=\left\{T_{i}^{c} \cap T_{j}^{d}\right\}
$$

If $I \neq \emptyset$ then $R_{i}$ and $R_{j}$ are merged. That is, in case there are some triangles created by region $R_{i}$ that are deleted by region $R_{j}$, then both regions must be merged. In the example of figure 3, first region $R_{i}$ was created (figures 3(a), 3(b)), then a new region $R_{j}$ is created from the collapse of edge between faces $n$ and $h$ (figures 3(c), 3(d)). As the intersection between $T_{i}^{c}$ and $T_{j}^{d}$ is $\{m, n\}$, both regions must be merged. We define the merge function as follows:

$$
R_{k}=\operatorname{merge}\left(R_{i}, R_{j}\right)=\left\{\left\{T_{i}^{d} \cup T_{j}^{d}\right\}-I,\left\{T_{i}^{c} \cup T_{j}^{c}\right\}-I\right\}
$$

Figures 3(e) and 3(f) show the new $T_{k}^{d}$ and $T_{k}^{c}$ sets of region $R_{k}$ as a result of the merge function. And the simplification error of the resulting region is the sum of the simplification errors of both regions. When a region reaches $E$, it is not further considered for the merging process. These completed regions are stored in a result list so that, at the end of the process, the list contains the sequence of regions in the completion order. Finally, for each region $R_{i}$ in the result list, a LoD of the multiresolution representation is established, that is, $M_{i}-T_{i}^{d}+T_{i}^{c}=M_{i-1}$. Figure 4 outlines the data structures and the algorithm used for this process.

## 5. Construction

One of the characteristics that a multiresolution model for real-time applications should fulfil is efficient information processing [RLB*02]. That is, if the multiresolution representation stores $n$ different LoDs, then the information should be organized in such a way that, during execution, when any of the $n$ LoDs is requested, the retrieval algorithm should extract data as fast as possible.

```
class Region {
    vector<int> DeletedFaces
    vector<int> CreatedFaces
    double error
}
list <Region> result
list <Region> temp
for each edge collapse e {
    construct new region }R(e
    for each region }\mp@subsup{R}{i}{}\mathrm{ in temp {
        if (\operatorname{check ( }\mp@subsup{R}{i}{},R)\not=\emptyset) {
            R= merge ( }\mp@subsup{R}{i}{},R
            R.error }=R.error + R,.error
            temp.remove ( }\mp@subsup{R}{i}{}\mathrm{ )
        }
    }
    if (R.error>E)
        result.push_back(R)
    else
        temp.push_back (R)
}
```

Figure 4: Algorithm for creating regions. The result list finally contains the sequence of regions in order of creation.

Let $M=M_{n}$ be the original mesh. Let $\left\{R_{n}, R_{n-1}, \ldots, R_{1}\right\}$, where $R_{i}=\left\{T_{i}^{d}, T_{i}^{c}\right\}$, the regions obtained in the simplification process. We construct each LoD as follows:

$$
\begin{aligned}
M_{n}-T_{n}^{d}+T_{n}^{c} & =M_{n-1} \\
& \cdots \\
M_{1}-T_{1}^{d}+T_{1}^{c} & =M_{0}
\end{aligned}
$$

Instead of storing every LoD independently, an incremental representation stores faces that compose the lower detailed LoD, $T_{0}$, plus the sequence of updates that allow to construct every LoD, which in our scheme are determined by the regions obtained in the simplification process.

$$
\begin{equation*}
M_{r}=\left\{\left\{T_{n}^{d}, T_{n}^{c}\right\},\left\{T_{n-1}^{d}, T_{n-1}^{c}\right\}, \ldots,\left\{T_{1}^{d}, T_{1}^{c}\right\}, T_{0}\right\} \tag{1}
\end{equation*}
$$

This representation would store many faces twice because a face created at $T_{i}^{c}$ most probably is deleted in some $T_{j}^{d}$ with $i>j$, but it can also remain until the lower detailed LoD (it belongs to $T_{0}$ ). Therefore, we can rewrite $M_{r}$ as a sequence of removed faces plus $T_{0}$, that is:

$$
\begin{equation*}
M_{r}=\left\{T_{n}^{d}, T_{n-1}^{d}, \ldots, T_{1}^{d}, T_{0}\right\} \tag{2}
\end{equation*}
$$

Each face is assigned a label which identifies the LoD where the face appears for the first time in the sequence $\left\{M_{n}, T_{n}^{c}, \ldots, T_{1}^{c}\right\}$. Faces that belong to $M_{n}$ are assigned the label $n$, those that belong to $T_{n}^{c}$ are assigned the label $n-1$, and so on, until those that belong to $T_{1}^{c}$, which are assigned the label zero. Consequently, each set $T_{i}^{d}$ probably contains triangles with different labels. Then, faces in each set $T_{i}^{d}$ are clustered depending on their label, so that the label is assigned to each cluster. Therefore, each set $T_{i}^{d}$ is composed of a set of labelled clusters, $T_{i_{l}}^{d}$. For a required LoD, $M_{i}$, a $T_{j_{l}}^{d}$, with $i \geq j$, belongs to $M_{i}$ only when $l \geq i$. Finally, in order to take advantage of graphics hardware features, each labelled cluster is processed to obtain triangle strips.

The basic data structure (figure 5) is composed of two vectors: one contains the vertices, and the other contains the sets of deleted faces plus $T_{0}$ (equation 2), that is, one set per LoD. Each set consists of an ordered vector of clusters. The number of clusters in each set $T_{i}^{d}$ depends on the number of labels in the set. Each cluster is composed of its label $l$ plus the set $T_{i_{l}}^{d}$ previously stripified. To speed up the recovery process, the clusters are stored in decreasing order of $l$.

```
class Vertex {float x, y, z}
class TriangleStrip {vector <uint> indices}
class Cluster {
    vector<TriangleStrip> triangleStrips
    int label
}
class Set {vector<Cluster> clusters}
class Mr {
    vector<Vertex> vertices
    vector<Set> sets
}
```

Figure 5: Basic data structure to represent $M_{r}$.

## 6. Uniform LoD recovery

The algorithm to retrieve a given LoD is very simple. Let $M_{i}, n \geq i \geq 0$, the required LoD to be rendered. As sets $T_{j}^{d}$, with $n \geq j>i$, are the sets of previously deleted faces, none of these faces belongs to $M_{i}$. So, the algorithm traverses the data structure starting from $T_{i}^{d}$ until $T_{0}$. Then, for each set $T_{j}^{d}, i \geq j \geq 0$, the clusters, $T_{j_{l}}^{d}$, such that $l \geq i$, belong to $M_{i}$. Let us remark that the algorithm does not need to check every triangle, only the label of each cluster. Also the ordering of clusters allows to speed up the process. The algorithm is shown in figure 6.

```
int First= sets.length()-i-1
int Last= sets.length()
for (j= First; j< Last; j++) {
    int nClusters= sets[j].clusters.length()
    int k= 0
    while ((k< nClusters) and
                (sets[j].clusters[k].label \geq i)) {
        Draw(sets[j].clusters[k])
        k= k +1
    }
}
```

Figure 6: Algorithm to render a uniform $\operatorname{LoD} M_{i}, n \geq i \geq 0$.


Figure 7: Number of clusters per set $T_{i}^{d}, 0 \leq i \leq 140$.

In order to analyze the efficiency of this algorithm, we must analyze the number of clusters per set, the total number of clusters in $M_{r}$, and the total number of clusters that form a given $\mathrm{LoD}, M_{i}$. First, let us analyze how many clusters, $T_{i_{l}}^{d}$, are there for each set $T_{i}^{d}$. Let $L_{i}$ be the number of clusters in the set $T_{i}^{d}$. Theoretically, the upper bound of $L_{i}$ is $n+$ $1-i$. The experimental measurements (figure 7) show that $L_{i}$ does not depend on the size of the object as expected. The variation of $L_{i}$ is very similar for all the objects and, in all cases, the values of $L_{i}$ are much less than the theoretical upper bound.

Second, the sum of every $L_{i}$ is bounded by $n^{2} / 2+n / 2$. Consequently, the maximum theoretical number of stored clusters is independent of the size of the object. Experimental results show that the total number of clusters is around $10 n$ for the tested objects (see table 1).

Third, let us analyze how many clusters form one LoD, $M_{i}$. The theoretical upper bound of the number of clusters that form $M_{i}$ is $i *(n+1-i)$, where $n \geq i \geq 0$. As $n$ depends on the desired frame rate, this amount is independent of the size of the object. Figure 8 shows the number of clusters per LoD of several objects, and the theoretical upper bound. The theoretical and experimental values coincide approximately in $M_{n}$ and $M_{0}$, but they are very different for the rest of LoDs. We can also observe that, on average, the maximum number of clusters that form a LoD is around $5 n$, which is very far


Figure 8: Number of clusters for each $M_{i}, 0 \leq i \leq 140$.
from the theoretical curve. Although the maximum differences between the objects are located around intermediate LoDs, they are almost insignificant, specially if we take into account, for example, that the dragon model has near 2 times the number of polygons of the armadillo model.

```
drawCluster (int i, vector<int> Labels) {
    int nClusters= sets[i].clusters.length()
    for (j= 0; j< nClusters; j++)
        if (sets[i].clusters[j].label }\not\in\mathrm{ Labels)
            Draw(sets[i].clusters[j])
}
checkDependence (int i, vector<int> Labels) {
    int nClusters= sets[i].clusters.length()
    int k= 0
    while ((k< nClusters) and
                    (sets[i].clusters[j].label & Labels))
            k= k + 1
    return (k< nClusters)
}
main () {
    vector<int> Labels= \emptyset
    int First= 0
    int Last= sets.length()
    for (i= First; i< Last; i++) {
        bool pDetail= preserveDetail(i)
        bool dependence= checkDependence(i,Labels)
        if ((pDetail is True) or
                (dependence is True)) {
            drawCluster(i,Labels)
            Labels.pushback (n-1-i)
        }
    }
}
```

Figure 9: Algorithm to render a view-dependent LoD.

## 7. View-Dependent LoD recovery

The aim is to represent an object with several levels of detail coexisting along the surface. The decision about which
areas of the surface should be visualized in high or low detail depends exclusively on the criterion or set of critera required by the application, for example, local illumination, view frustum, silhouette and so on [Hop97].
The algorithm to retrieve a view-dependent LoD evaluates each set $T_{i}^{d}$ in order to decide whether its faces must be preserved or removed in the required LoD, by using a function that evaluates the criterion defined by the application. Therefore, the algorithm traverses the whole data structure starting from $T_{n}^{d}$ until $T_{0}$ (see equation 2). Let us note that, although this might seem a drawback, the vector length is independent of the object complexity, and the experimental results show it is around 10 n . If the evaluation function finds that some face in $T_{i}^{d}$ must be preserved, all faces in $T_{i}^{d}$ are preserved. That is, the simplification coded by $T_{i}^{d}$ is not performed. Therefore, none of the faces in $T_{i}^{c}$ can belong to the required LoD, because these faces are created only if the simplification coded by $T_{i}^{d}$ is performed. Then, if one set $T_{j}^{d}, j \neq i$, contains any of the faces in $T_{i}^{c}$, the simplification coded by $T_{j}^{d}$ can not be performed, as it depends on the simplification coded by $T_{i}^{d}$. Consequently, faces in $T_{j}^{d}$ not belonging to $T_{i}^{c}$ are preserved too. Actually, only sets $T_{j}^{d}$ with $i>j$ can suffer this dependence and, in case the dependence ocurrs, $T_{j}^{d}$ also generates new dependences with subsequent sets. As faces in $T_{j}^{d}$ have been clustered and labelled, and each label indicates the LoD of creation, none of the faces in cluster $T_{j_{l}}^{d}$ with $l \neq i-1$ belong to $T_{i}^{c}$, only faces in $T_{j_{i-1}}^{d}$ do. Therefore, for checking dependence between $T_{i}^{d}$ and $T_{j}^{d}, i>j$, it is enough to check if any of the labels in $T_{j}^{d}$ is equal to $i-1$. The algorithm is shown in figure 9. In summary, each set $T_{i}^{d}, n \geq i \geq 0$, is preserved in any of these cases:

- The evaluation function decides to preserve detail.
- There exists a dependence: there is one set $T_{k}^{d}, k>i$, that was preserved and there is a cluster $T_{i_{l}}^{d}$, such that $l=k-1$.
This process uses the same data structure shown in figure 5. However, the algorithm uses a vector of labels to store those which produce dependencies in the recovered LoD. The length of this vector is $n+1$ in the worst case, where $n$ is a constant multiple of the frame rate. The algorithm traverses the sets vector entirely. As the length of this vector is small (50, 100 and 140 in our experiments), the traversal is done very fast. Probably, the most expensive operation in the algorithm is the evaluation function for preserving detail. So, the interest of the algorithm depends on whether this cost is lower than the cost in rendering the most detailed LoD using the algorithm shown in figure 6 .


## 8. Results

The experiments were carried out on a personal computer with Linux operating system, Pentium D CPU and NVIDIA GeForce $7600-\mathrm{GT}$ GPU. The model was coded in C++ and

|  | Geometry. $M_{n}$ |  | $M_{r}, 140$ LoDs |  | $M_{r}, 100$ LoDs |  | $M_{r}, 50$ LoDs |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
|  | \#Vertices | \#Triangles | \#Indices | \#Strips | \#Indices | \#Strips | \#Indices | \#Strips |
| Armadillo | 172,974 | 345,944 | $1,021,755$ | 1,330 | $1,010,339$ | 951 | 975,797 | 436 |
| Dragon | 359,173 | 715,933 | $2,069,974$ | 1.577 | $2,053,243$ | 1,171 | $2,005,531$ | 535 |

Table 1: Three iDLoD representations of Armadillo and Dragon objects, with 140 LoDs, 100 LoDs and 50 LoDs respectively.

|  | $M_{100}$ |  | $M_{75}$ |  | $M_{50}$ |  | $M_{25}$ |  | $M_{0}$ |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
|  | \#Vertices | \#Indices | \#Vertices | \#ndices | \#Vertices | \#Indices | \#Vertices | \#Indices | \#Vertices | \#Indices |
| Armadillo | 172,974 | 549,196 | 129,956 | 421,305 | 86,939 | 285,610 | 43,921 | 145,241 | 904 | 3,042 |
| Dragon | 359,173 | $1,030,325$ | 270,463 | 841,201 | 181,770 | 580,827 | 93,072 | 301,434 | 3,710 | 11,250 |

Table 2: DLoD representations of Armadillo and Dragon objects with 5 LoDs, each one reducing the geometry in $25 \%$.
the OpenGL graphics library was used. In our experiments we enabled one light source and the size of the viewport is $1024 \times 768$ pixels. For each vertex we send its coordinates and its normal, and we scale the object to the bigger size inside of the frustum. We used the GL_TIME_ELAPSED extension, which provides a query mechanism to determine the amount of time used for completing a set of GL tasks without stalling the rendering pipeline. We used the NvTriStrip library for vertex cache aware stripification of geometry. Experiments show that better performance is obtained when triangle strips are stitched together using degenerate triangles. Therefore, triangle strips have been stiched in the off-line process for both DLoD and iDLoD.

Table 1 shows the characteristics of the polygonal models used in the experiments as well as the number of strip indices. The number of clusters in $M_{r}$ is equal to the number of strips since they are stiched for each cluster. The experiments carried out aim to obtain the performance of the proposed model, iDLoD, and even more important, to compare with the performance of a DLoD that consists of 5 LoDs , denoted as reference LoDs, each one reducing the geometry in $25 \%$. Table 2 show the characterisitcs of the DLoD constructed for each object.

We used two different implementations of the iDLoD model. First, a unique multiresolution representation with the $n$ LoDs, denoted as iDLoD. Second, a sequence of four multiresolution representations, denoted as iDLoD-b, such that each one provides LoDs between two consecutive reference LoDs. Therefore, each multiresolution representation in the iDLoD-b model provides a quarter of the $n$ LoDs. Figure 10 shows the results of the three considered models with two different objects. All of the plots show similar behaviour. As expected, the DLoD produces the highest frame rate. The iDLoD almost reaches DLoD performance: the models with 50 LoDs approach DLoD performance more than the ones with 140 LoDs . The performance of the iDLoD-b implementation goes through the reference values of DLoDs. This is because this implementation provides LoDs formed by a quarter of the clusters at most,
which produces the increment of performance respect to the iDLoD implementation.

## 9. Conclusions

Since J. Clark presented it in 1976, DLoD has been the multiresolution framework mostly used. During the last years, many well-known works have been proposed as an alternative. Mainly, the efforts have been directed to improve CLoD and View-Dependent techniques. However, in the field of real time applications, software developers still prefer the DLoD framework. In this paper we have presented a model based on coding a discrete number of levels of detail, with more LoDs coded than is usual in DLoD, and with an incremental representation, which is often used in CLoD. The difference between two consecutive LoDs is forced to be a surface patch so that it can be optimized off-line for rendering. Experimental results show that this model obtains a performance similar to DLoD by providing optimized LoDs for efficient visualization, while the popping effect is imperceptible. In addition, view-dependent LoDs can be retrieved in spite of the low granularity.

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Figure 10: Results for the Armadillo and Dragon objects with different number of LoDs.
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